

Local Planning

$$A_{s_0} = \operatorname{argmax}_a (\hat{T}_{s_0}^H v)(s_0, a)$$

What does this buy us??

Localizing computation to s_0

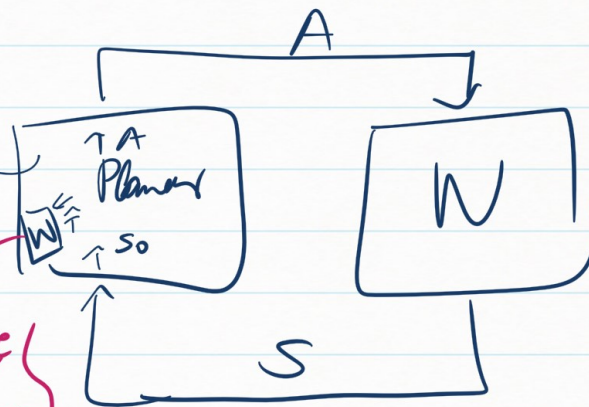
Goal? | $A = a_1$

① Little compute

② Input: $\delta > 0$

δ -optimality

π



Planner $\rightarrow \pi(a|s_0) = \mathbb{P}_{s_0}(A_{s_0}=a)$

$$v^{\pi} \geq v^* - \delta 1$$

per-state compute cost

$$O((AV\delta)^H) \dots$$



$$(s, a) \mapsto \underbrace{S'_1(s, a), \dots, S'_m(s, a)}_{\sim P_a(s)} \leftarrow C(s, a)$$

independencia

"call simulator"

$$C(s) = (C(s, a))_{a \in A} / \cdot = \bigcup_{a \in A} C(s, a)$$

$$\hat{T}: \mathbb{R}^{S \times A} \rightarrow \mathbb{R}^{S \times A}$$

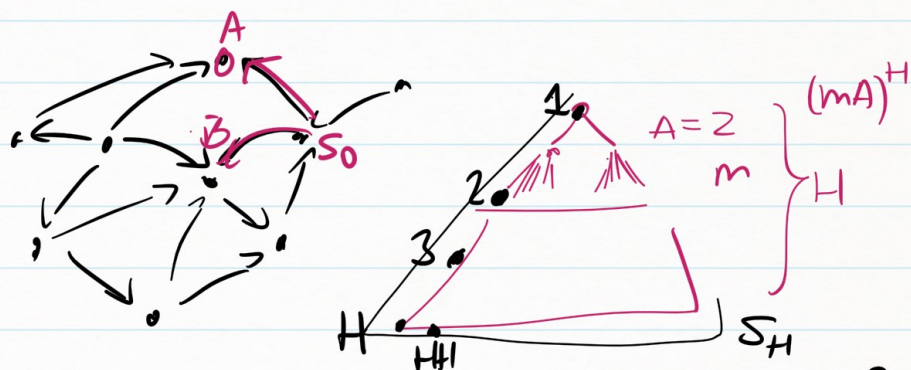
$(Mq)(s')$

$$(\hat{T}q)(s, a) = r_a(s) + \gamma \frac{1}{m} \sum_{s' \in C(s, a)} \max_{a'} q(s', a')$$

$$(\hat{T}q)(s, a) = r_a(s) + \gamma \langle P_a(s), Mq \rangle$$

$$G = (S, \mathcal{E})$$

$$\mathcal{E} = \{(s, s') \mid s' \in C(s)\}$$



$$S_h = \{s \in S \mid \text{dist}(s_0, s) \leq h\}$$

$$h=0 \quad S_0 = \{s_0\}$$

$$S_1 = \{s_0, A, B\}$$

\vdots

$$\delta_h = \|\hat{T}^h 0 - q^*\|_{S_{H-h}}$$

$$\delta_h \leq \gamma \delta_h + \underbrace{\|\hat{T}q^* - Tq^*\|_{S_H}}_{\leq \frac{\epsilon}{1-\gamma}} \text{ "whp"}$$

$$\delta_0 \leq \frac{1}{1-\gamma}$$

$$\delta_h \leq \gamma \delta_{h-1} + \frac{\varepsilon'}{1-\gamma} \quad , \quad h \geq 1$$

$$\delta_0 \leq \frac{1}{1-\gamma}$$

$$\delta_1 \leq \frac{\gamma + \varepsilon'}{1-\gamma}$$

$$\delta_2 \leq \frac{\gamma(\gamma + \varepsilon') + \varepsilon'}{1-\gamma} = \frac{\gamma^2 + \varepsilon'(1+\gamma)}{1-\gamma}$$

$$\vdots$$

$$\delta_H \leq \frac{\gamma^H + \varepsilon'(1 + \gamma + \dots + \gamma^{H-1})}{1-\gamma}$$

$$\leq \left(\gamma^H + \frac{\varepsilon'}{1-\gamma} \right) \frac{1}{1-\gamma} \leq \delta \quad \boxed{Mq^* = v^*}$$

$$\boxed{\varepsilon' = ?}$$

$$\max_{S \in \mathcal{S}_H} \max_a \gamma \left(\frac{1}{m} \sum_{j=1}^m v^*(S_j^i(s_a)) - \langle P_a, \theta \rangle \right) \rightarrow \underbrace{\|\hat{T}q^* - Tq^*\|_{\mathcal{S}_H}}_{\leq \frac{1}{1-\gamma}}$$

Hoeffding's inequality

Lemma: $0 \leq X_i \leq 1$ i.i.d.
 $i = 1, \dots, n$

$$(C) \left\{ \begin{array}{l} \forall 0 \leq \zeta < 1 \quad \text{wp at least } 1-\zeta, \\ \left| \frac{1}{n} \sum_{i=1}^n X_i - \mathbb{E} X_1 \right| \leq \sqrt{\frac{\log(2/\zeta)}{2n}} \end{array} \right.$$

Lemma: $0 \leq X_i \leq 1, i = 1, \dots, n$
random element U

(U, X_1, \dots, X_n) jointly distributed.

X_1, \dots, X_n i.i.d. given U

$$[P(X_i \in A_i, i=1 \dots n | U) = \prod_{i=1}^n P(X_i \in A_i | U)]$$

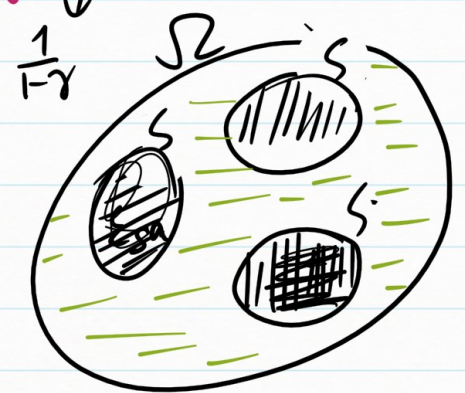
$$\forall 0 \leq \zeta < 1 \quad \text{wp at least } 1-\zeta$$

$$\left| \frac{1}{n} \sum_{i=1}^n X_i - \mathbb{E}[X_1 | U] \right| \leq \sqrt{\frac{\log(2/\zeta)}{2n}}$$

$$\|\hat{T}_{q^*} - T_{q^*}\|_{S_H} = \max_{S \in S_H} \max_{a \in A} \left| \frac{1}{m} \sum_{j=1}^m v^*(S_j(s, a)) - \langle P_a(s), v^* \rangle \right|$$

Union bounds!

(outside of $\mathcal{E}_{S,a}$)
 $\forall (s, a) \quad |\Delta(s, a)| \leq \frac{1}{1-\gamma} \sqrt{\frac{\log(2/S)}{m}}$
 $\mathbb{P}(\mathcal{E}_{S,a}) \leq \dots$

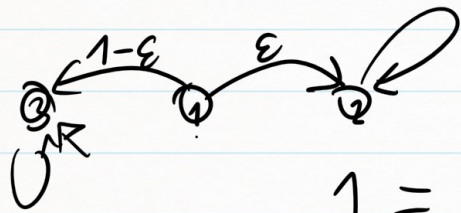


Outside of $\bigcup_{S,a} \mathcal{E}_{S,a} \Rightarrow \max_{(s,a)} |\Delta(s,a)| \leq \frac{1}{1-\gamma} \sqrt{\frac{\log(2/S)}{2m}}$

$$\mathbb{P}\left(\left(\bigcup_{S,a} \mathcal{E}_{S,a}\right)^c\right) = 1 - \mathbb{P}\left(\bigcup_{S,a} \mathcal{E}_{S,a}\right) \geq 1 - \sum_{S,a} \mathbb{P}(\mathcal{E}_{S,a}) \geq 1 - SA\gamma$$

w.p. $1 - \delta_{\text{failure}}$
 $\delta_{\text{failure}} = SA\gamma$
 $\gamma = \frac{\delta_{\text{failure}}}{SA}$
 $|S_H| \leq (mA)^{H+1}$ (random)
 take union bound over these?!

$\frac{\epsilon'}{1-\gamma} \leq \frac{1}{1-\gamma} \sqrt{\frac{\log(2SA)}{\delta_{\text{failure}}}} \cdot \frac{1}{2m}$
 $H = \frac{1}{(1-\gamma)^2}$
 $\frac{2\delta_{\text{failure}}}{(1-\gamma)^2} = \frac{\gamma}{3} \leftarrow \text{target subset.}$
 $m \approx H \log(Am)$
 $m \approx H \log(A)$
 $m = \Omega(\log(S))$



$$S' \sim P_A$$

(S) random

$$1 = \mathbb{P}(S'=2 | S_1=2) =$$

$$\mathbb{P}(S'=2)$$

$\approx \varepsilon$

$$S \doteq S'$$

$\{\hat{S}_1, \dots, \hat{S}_n\}$
 $\forall A \leq i \leq n$

$$S_H = \{S_1, \dots, S_n\}$$

$$n = (Am)^{t+i} - 1 / A - 1$$

S_1, \dots, S_n depth-first

encounters of states in S_H

Lemma:

$\forall i \in [n]$ wp $1 - \delta$
 $\forall a \in [A]$

$$|\Delta(S_i, a)| \leq \frac{1}{\sqrt{r}} \sqrt{\frac{\log(2/\delta)}{2m}}$$

Proof: $\forall i, a$

$$\rightarrow J_i = \min \{1 \leq k \leq i \mid S_k = S_i\}$$

$$1 \leq J_i \leq i$$

$$\Delta(S_i, a) = \frac{1}{m} \sum_{s' \in C(S_i, a)} v^a(s') - \underbrace{P_a(S_i, v^a)}_{S_i}$$

$$\Delta(S_i, a) = \Delta(S_{J_i}, a)$$

$C(S_{J_i}, a) \perp S_{J_i}$ gia S_{J_i}
 $\{C(S_1), \dots, C(S_{J_i-1})\}$

+ union bound over $i \in [n]$

+ cranking together

Thm: $\text{Cost} = O((mA)^H)$

$$H \approx H_{\gamma, (1-\gamma)^2 \delta}$$

$$m \approx H \log(A)$$

local plugging error π :

$$v^{\pi} \geq v^* - \delta \mathbb{1}$$

Kearns, Mansour, Singh 2018

